## Package ‘deducorrect’

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deducorrect-package Deductive correction methods for sign, rounding and typing errors

## Description

Deductive correction methods for sign, rounding and typing errors It also contains functionality to check if a matrix of linear restrictions is totally unimodular. See also

## Details

- correctRounding
- correctTypos
- deducorrect-object and status for output specification
- isTotallyUnimodular
- deduImpute
- solSpace, imputess
- deductiveZeros
- deductiveLevels


## Examples

```
require(editrules)
# some data
dat <- data.frame(
    x = c( 3, 14, 15),
    y = c(13, -4, 5),
    z = c(10, 10,-10))
dat
# ... which has to obey
E <- editmatrix("z == x-y")
```

```
    # All signs may be flipped, no swaps.
    correctSigns(E, dat)
correctTypos(E, dat)
correctRounding(E, dat)
```

correctionRules Rules for deterministic correction

## Description

Rules for deterministic correction

## Usage

```
    correctionRules(x, strict = TRUE, allowed = getOption("allowedSymbols"),
        ...)
    \#\# S3 method for class 'character'
    correctionRules(x, strict = TRUE,
        allowed = getOption("allowedSymbols"), file = TRUE, ...)
    \#\# S3 method for class 'expression'
    correctionRules(x, strict = TRUE,
        allowed = getOption("allowedSymbols"), ...)
    \#\# S3 method for class 'correctionRules'
    print(x, ...)
    \#\# S3 method for class 'correctionRules'
    as.character (x, oneliner = FALSE, ...)
    \#\# S3 method for class 'correctionRules'
    getVars(E, ...)
```


## Arguments

X
strict
... Currently unused.
oneliner Coerce to oneliner
E
allowed A character vector of allowed symbols
file If file=TRUE, $x$ is treated as a filename from which the rules are read.
character or expression vector.
If TRUE an error is thrown if any forbidden symbol is used (see details).
object of class correctionRules

## Value

correctionRules returns an object of class correctionRules
getVars returns a character vector of variable names.

## Details

Data editing processes are rarely completely governed by in-record consistency rules. Many ad$h o c$ rules are commonly used to impute empty or erroneous values. Such rules are often applied manually or hidden in source code. This function, together with correctWithRules allows for easy definition and execution of simle deterministic replacement rules.
These functions are ment to support very simple rules, such as if variable $x$ is missing, then set it to zero. Such actions usually basically model-free corrections stemming from subject-matter knowledge. Given the nature of such rules, the type of rules are by default limited to R-statements containing conditionals (if-else), arithmetic and logical operators, and brackets and assignment operators. see getOption('allowedSymbols') for a complete list.
If you cannot execute your 'simple' corrections with just these functions, we strongly recommend to write a separate imputation or correction routine. However, it's a free world, so you may alter the list of allowed symbols as you wish.

## Note

getVars is overloaded from the editrules package.

## See Also

correctWithRules

## Description

This algorithm tries to detect and repair records that violate linear (in)equality constraints by correcting possible rounding errors as described by Scholtus(2008). Typically data is constrainted by $R x=a$ and $Q x \geq b$.

## Usage

```
correctRounding(E, dat, ...)
## S3 method for class 'editset'
correctRounding(E, dat, ...)
## S3 method for class 'editmatrix'
correctRounding(E, dat, fixate = NULL, delta = 2,
    K = 10, round = TRUE, assumeUnimodularity = FALSE, ...)
```


## Arguments

E editmatrix or editset as generated by the editrules package.
dat data.frame with the data to be corrected
... arguments to be passed to other methods.
fixate character with variable names that should not be changed.
delta tolerance on checking for rounding error
K number of trials per record. See details
round should the solution be rounded, default TRUE
assumeUnimodularity
If FALSE, a test is performed before corrections are computed (expensive).

## Details

The algorithm first finds violated constraints $\left|r_{i}^{\prime} x-a_{i}\right|>0$, and selects edits that may be due to a rounding error $0<\left|r_{i}^{\prime} x-a_{i}\right| \leq \delta$. The algorithm then makes a correction suggestion where the errors are attributed to randomly selected variables under the lineair equality constraints. It checks if the suggested correction does not violate the inequality matrix $Q$. If it does, it will try to generate a different solution up till K times.

## Value

A deducorrrect object.

## References

Scholtus S (2008). Algorithms for correcting some obvious inconsistencies and rounding errors in business survey data. Technical Report 08015, Statistics Netherlands.

## See Also

deducorrect-object status

## Examples

```
E <- editmatrix(expression(
    x1 + x2 == x3,
    x2 == x4,
    x5 + x6 + x7 == x8,
    x3 + x8 == x9,
    x9 - x10 == x11
    )
)
dat <- data.frame( x1=12
            , x2=4
            , x3=15
            , x4=4
```

```
, x5=3
, x6=1
, x7=8
, x8=11
, x9=27
, x10=41
x11=-13
)
sol <- correctRounding(E, dat)
# example with editset
for ( d in dir("../pkg/R/",full.names=TRUE) ) dmp <- source(d)
E <- editmatrix(expression(
    x + y == z,
    x >= 0,
    y >= 0,
    z >= 0,
    if (x>0) y > 0
    ))
dat <- data.frame(
    x = 1,
    y = 0,
    z = 1)
# solutions causing new violations of conditional rules are rejected
sol <- correctRounding(E,dat)
# An example with editset
E <- editset(expression(
    x + y == z,
    x >= 0,
    y>0,
    y<2,
    z> 1,
    z<3,
    A %in% c('a','b'),
    B %in% c('c','d'),
    if ( A == 'a' ) B == 'b',
    if ( B == 'b' ) x < 1
))
dat <- data.frame(
    x = 0,
    y = 1,
    z = 2,
    A = 'a',
    B = 'b'
)
correctRounding(E,dat)
```


## Description

Correct sign errors and value interchanges in data records.

## Usage

```
correctSigns(E, dat, ...)
## S3 method for class 'editset'
correctSigns(E, dat, ...)
## S3 method for class 'editmatrix'
correctSigns(E, dat, flip = getVars(E), swap = list(),
    maxActions = length(flip) + length(swap), maxCombinations = 1e+05,
    eps = sqrt(.Machine$double.eps), weight = rep(1, length(flip) +
    length(swap)), fixate = NA, ...)
```


## Arguments

E
dat data. frame, the records to correct.
... arguments to be passed to other methods.
flip A character vector of variable names who's values may be sign-flipped
swap A list of character 2-vectors of variable combinations who's values may be swapped
maxActions The maximum number of flips and swaps that may be performed maxCombinations

The number of possible flip/swap combinations in each step of the algorithm is choose ( $n, k$ ), with $n$ the number of flips+swaps, and $k$ the number of actions taken in that step. If choose ( $n, k$ ) exceeds maxCombinations, the algorithm returns a record uncorrected.
eps Tolerance to check equalities against. Use this to account for sign errors masked by rounding errors.
weight weight vector. Weights can be assigned either to actions (flips and swap) or to variables. If length (weight) $==1$ length ( $f$ lip) + length ( $s w a p$ ), weights are assiged to actions, if length (weight) $==n \operatorname{col}(E)$, weights are assigned to variables. In the first case, the first length\{flip\} weights correspond to flips, the rest to swaps. A warning is issued in the second case when the weight vector is not named. See the examples for more details.
fixate a character vector with names of variables whos values may not be changed

## Details

This algorithm tries to correct records violating linear equalities by sign flipping and/or value interchanges. Linear inequalities are taken into account when judging possible solutions. If one or more inequality restriction is violated, the solution is rejected. It is important to note that the status of a record has the following meaning:

$$
\begin{array}{ll}
\text { valid } & \text { The record obeys all equality constraints on entry. No error correction is performed. } \\
\text { It may therefore still contain inequality errors. } \\
\text { corrected } & \text { Equality errors were found, and all of them are solved without violating inequalities. } \\
\text { partial } & \text { Does not occur } \\
\text { invalid } & \begin{array}{l}
\text { The record contains equality violations which could not be solved with this algorithm } \\
\text { NA }
\end{array} \\
\text { record could not be checked. It contained missings. }
\end{array}
$$

The algorithm applies all combinations of (user-allowed) flip- and swap combinations to find a solution, and minimizes the number of actions (flips+swaps) that have to be taken to correct a record. When multiple solutions are found, the solution of minimal weight is chosen. The user may provide a weight vector with weights for every flip and every swap, or a named weight vector with a weight for every variable. If the weights do not single out a solution, the first one found is chosen.
If arguments flip or swap contain a variable not in $E$, these variables will be ignored by the algorithm.

## Value

a deducorrect-object. The status slot has the following columns for every records in dat.

| status | a status factor, showing the status of the treated record. |
| :--- | :--- |
| degeneracy | the number of solutions found, after applying the weight |
| weight | the weight of the chosen solution |
| nflip | the number of applied sign flips |
| nswap | the number of applied value interchanges |

## References

Scholtus S (2008). Algorithms for correcting some obvious inconsistencies and rounding errors in business survey data. Technical Report 08015, Netherlands.

## See Also

deducorrect-object

## Examples

```
# some data
dat <- data.frame(
    x = c( 3,14,15, 1, 17,12.3),
    y = c(13,-4, 5, 2, 7, -2.1),
    z = c(10,10,-10,NA,10,10 ))
```

```
# ... which has to obey
E <- editmatrix(c("z == x-y"))
# All signs may be flipped, no swaps.
correctSigns(E, dat)
# Allow for rounding errors
correctSigns(E, dat, eps=2)
# Limit the number of combinations that may be tested
correctSigns(E, dat, maxCombinations=2)
# fix z, flip everything else
correctSigns(E, dat,fixate="z")
# the same result is achieved with
correctSigns(E, dat, flip=c("x","y"))
# make x and y swappable, allow no flips
correctSigns(E, dat, flip=c(), swap=list(c("x","y")))
# make x and y swappable, swap a counts as one flip
correctSigns(E, dat, flip="z", swap=list(c("x","y")))
# same, but now, swapping is preferred (has lower weight)
correctSigns(E, dat, flip="z", swap=list(c("x","y")), weight=c(2,1))
# same, but now becayse x any y carry lower weight. Also allow for rounding errors
correctSigns(E, dat, flip="z", swap=list(c("x","y")), eps=2, weight=c(x=1, y=1, z=3))
# demand that solution has y>0
E <- editmatrix(c("z==x-y", "y>0"))
correctSigns(E,dat)
# demand that solution has y>0, taking acount of roundings in equalities
correctSigns(E,dat,eps=2)
# example with editset
E <- editset(expression(
    x + y == z,
    x >= 0,
    y>0,
    y<2,
    z>1,
    z<3,
    A %in% c('a','b'),
    B %in% c('c','d'),
    if ( A == 'a' ) B == 'b',
    if ( B == 'b' ) x < 1
))
x <- data.frame(
```

```
    x = -1,
    y = 1,
    z = 2,
    A = 'a',
    B = 'b'
)
correctSigns(E,x)
```

correctTypos Correct records under linear restrictions using typographical error suggestions

## Description

This algorithm tries to detect and repair records that violate linear equality constraints by correcting simple typo's as described in Scholtus (2009). The implemention of the detection of typing errors differs in that it uses the restricted Damerau-Levensthein distance. Furthermore it solves a broader class of problems: the original paper describes the class of equalities: $E x=0$ (balance edits) and this implementation allows for $E x=a$.

## Usage

correctTypos(E, dat, ...)
\#\# S3 method for class 'editset'
correctTypos(E, dat, ...)
\#\# S3 method for class 'editmatrix'
correctTypos(E, dat, fixate $=$ NULL, cost $=c(1,1,1$,
1), eps = sqrt(.Machine\$double.eps), maxdist $=1, \ldots$ )

## Arguments

E
dat data.frame with data to be corrected.
... arguments to be passed to other methods.
fixate character with variable names that should not be changed.
cost for a deletion, insertion, substition or transposition.
eps numeric, tolerance on edit check. Default value is sqrt(.Machine\$double.eps). Set to 2 to allow for rounding errors. Set this parameter to 0 for exact checking.
maxdist numeric, tolerance used in finding typographical corrections. Default value 1 allows for one error. Used in combination with cost.

## Details

For each row in dat the correction algorithm first detects if row x violates the equality constraints of E taking possible rounding errors into account. Mathematically: $\left|\sum_{i=1}^{n} E_{j i} x_{i}-a_{j}\right| \leq \varepsilon, \quad \forall j$
It then generates correction suggestions by deriving alternative values for variables only involved in the violated edits. The correction suggestions must be within a typographical edit distance (default $=1)$ to be selected. If there are more then 1 solutions possible the algorithm tries to derive a partial solution, otherwise the solution is applied to the data.
correctTypos returns an object of class deducorrect object describing the status of the record and the corrections that have been applied.
Inequalities in editmatrix $E$ will be ignored in this algorithm, so if this is the case, the corrected records are valid according to the equality restrictions, but may be incorrect for the given inequalities.

Please note that if the returned status of a record is "partial" the corrected record still is not valid. The partially corrected record will contain less errors and will violate less constraints. Also note that the status "valid" and "corrected" have to be interpreted in combination with eps. A common scenario is first to correct for typo's and then correct for rounding errors. This means that in the first step the algorithm should allow for typo's (e.g. eps==2). The returned "valid" record therefore may still contain rounding errors.

## Value

deducorrect object with corrected data.frame, applied corrections and status of the records.

## References

Scholtus S (2009). Automatic correction of simple typing errors in numerical data with balance edits. Discussion paper 09046, Statistics Netherlands, The Hague/Heerlen.
Damerau F (1964). A technique for computer detection and correction of spelling errors. Communications of the ACM, 7, issue 3
Levenshtein VI (1966). Binary codes capable of correcting deletions, insertions, and reversals. Soviet Physics Doklady 10: 707-10
A good description of the restricted DL-distance can be found on wikipedia: http://en.wikipedia.org/wiki/Damerau

## See Also

damerauLevenshteinDistance

## Examples

```
library(editrules)
# example from section 4 in Scholtus (2009)
E <- editmatrix( c("x1 + x2 == x3"
    ,"x2 == x4"
    ,"x5 + x6 + x7 == x8"
    ,"x3 + x8 == x9"
```

```
        ,"x9 - x10 == x11"
        )
)
dat <- read.csv(txt<-textConnection(
" , x1, x2 , x3 , x4 , x5 , x6, x7, x8 , x9 , x10 , x11
4 , 1452, 116, 1568, 116, 323, 76, 12, 411, 1979, 1842, 137
4.1, 1452, 116, 1568, 161, 323, 76, 12, 411, 1979, 1842, 137
4.2, 1452, 116, 1568, 161, 323, 76, 12, 411, 19979, 1842, 137
4.3, 1452, 116, 1568, 161, 0, 0, 0, 411, 19979, 1842, 137
4.4, 1452, 116, 1568, 161, 323, 76, 12, 0, 19979, 1842, 137"
))
close(txt)
(cor <- correctTypos(E,dat))
# example with editset
E <- editset(expression(
    x + y == z,
    x >= 0,
    y>0,
    y<2,
    z > 1,
    z< 3,
    A %in% c('a','b'),
    B %in% c('c','d'),
    if ( A == 'a' ) B == 'b',
    if ( B == 'b' ) x > 3
))
x <- data.frame(
    x = 10,
    y = 1,
    z = 2,
    A = 'a',
    B = 'b'
)
correctTypos(E,x)
```

correctWithRules Deterministic correction

## Description

Apply simple replacement rules to a data. frame.

## Usage

correctWithRules(rules, dat, strict = TRUE)

## Arguments

```
    rules object of class correctionRules
    dat data.frame
    strict If TRUE, an error is produced when the rules use variables other than in the
            data.frame.
```


## Value

list with altered data (\$corrected) and a list of alterations (\$corrections).

## Details

This function applies the the rules one by one to dat and logs their actions. Rules are excuted in order of occurrence in the correctionRules so order may matter for the final result. Rules are applied to one record at the time, so the use of statistical funtions such as mean is useless, and forbidden by default. See correctionRules for details on the type of rules that are possible.

## See Also

correctionRules

## Examples

```
## Some example data
dat <- data.frame(
    x = c(NA, 2, 0,-10),
    y = c(1, NA,NA, 6)
)
## a few rules
u <- correctionRules(expression(
    if ( is.na(x) ) x <- 0,
    if ( x == 0 && is.na(y) ) y <- 0,
    if ( is.na(y) ) y <- 1,
    if ( x < 0 ) y <- 0
))
correctWithRules(u,dat)
```


## Description

The restricted Damerau Levenshtein Distance between two strings is commonly used for checking typographical errors in strings. It takes the deletion and insertion of a character, a wrong character (substition) or the swapping (transposition) of two characters into account. By default these operations each account for distance 1.

## Usage

damerauLevenshteinDistance(sa, sb, w = c(1, 1, 1, 1))

## Arguments

sa character vector
sb character vector of equal length (sa)
w integer vector for cost of deletion, insertion, substitution and transposition.

## Value

integer vector with pairwise edit distances

## References

Damerau F (1964). A technique for computer detection and correction of spelling errors. Communications of the ACM, 7, issue 3

Levenshtein VI (1966). Binary codes capable of correcting deletions, insertions, and reversals. Soviet Physics Doklady 10: 707-10 Damerau Levenshtein Distance calculates the difference between two strings used for typographical errors (typo's)

## Description

deducorrect object

## Details

All correct* functions of the deducorrect package return an S3 object of class deducorrect. The "public slots" are

- corrected A copy of the input data.frame, with records corrected where possible.
- corrections A data.frame with the following colums:
- row Row number where correction was applied
- variable Variable name where correction was applied
- old Old value of adapted variable
- new New value of adapted variable
- status A data.frame with the same number of rows as corrected. It has at least a column called status. Further columns might be present, depending on the used correction function.
- generatedby The name of the function that called newdeducorrect to construct the object.
- timestamp The time and date when the object was generated, as returned by date.
- user The system's username of the user running R. Note that this may yield unexpected results when R accessed on a remote (web)server.
deductiveLevels Derive imputation values for categorical data


## Description

Deduce imputation values for categorical data. By substituting all known values and interatively eliminating the unknowns from the set of edits, unique imputation values are derived where possible.

## Usage

deductiveLevels(E, x, adapt = rep(FALSE, length(x)), checkFeasibility = TRUE)

## Arguments

E editarray
$x \quad$ a named character vector
adapt boolean vector indicating which variables may be adapted.
checkFeasibility
Test wether the assumed-correct values (observed and not designated by adapt) can lead to a consistent record.

## Details

Imputation values are derived for missing variables (NA) and for variables indicated by 'adapt'.

## Value

A named vector with imputation values for x

## References

T. De Waal, J. Pannekoek and S. Scholtus (2011) Handbook of statistical data editing Chpt 9.2.19.2.2

## Examples

```
# a simple example. We know the subject is pregnant. What is the gender?
E <- editarray(c(
    "gender \%in\% c('male','female')",
    "pregnant \%in\% c(TRUE,FALSE)",
    "if ( gender=='male') !pregnant"))
# a record with unknown gender
x <- c(gender=NA,pregnant=TRUE)
# find imputations
(s <- deductiveLevels(E,x))
# imputation can be done as follows:
x[names(s)] <- s
# Here's an example from De Waal et al (2011) (ex. 9.3)
E <- editarray(c(
    "x1 \%in\% letters[1:4]",
    "x2 \%in\% letters[1:3]",
    "x3 \%in\% letters[1:3]",
    "x4 \%in\% letters[1:2]",
    "if (x2 == 'c' & x3 != 'c' & x4 == 'a' ) FALSE",
    "if (x2 != 'a' & x4 == 'b') FALSE",
    "if (x1 != 'c' & x2 != 'b' & x3 != 'a') FALSE",
    "if (x1 == 'c' & x3 != 'a' & x4 == 'a' ) FALSE"
))
x<- c(x1='c',x2='b',x3=NA,x4=NA)
(s <- deductiveLevels(E,x))
x[names(s)] <- s
# another example, partial imputation
y <- c(x1=NA, x2=NA,x3=NA, x4='b')
(s <- deductiveLevels(E,y))
y[names(s)]<- s
```


## Description

Find out which variables can deductively be imputed with 0
Interface for deductiveZeros for objects of class editmatrix. This interface is robust for variables in $x$ not occuring in $E$.
Suppose $x$ is a record under linear constraints $A x=b$ and $x \geq 0$. In certain cases some missing values can be imputed uniquely with zeros. For example, in the case that $x_{1}+x_{2}=x_{3}$, if $x_{2}$ is missing and $x_{1}=x_{3} \geq 0$, then $x_{2}$ is uniquely determined to be 0 . This function returns a boolean vector indicating which of the missing values are uniquely determined to be zero.

## Usage

```
deductiveZeros(E, x, ...)
## S3 method for class 'editmatrix'
deductiveZeros(E, x, ...)
## S3 method for class 'matrix'
deductiveZeros(E, x, b, adapt = logical(length(x)),
    nonneg = rep(TRUE, length(x)), roundNearZeros = TRUE,
    tol = sqrt(.Machine$double.eps), ...)
```


## Arguments

E
$x \quad$ named numeric vector. Naming is optional if E is an equality constraint matrix.
... extra parameters to pass to deductiveZeros, matrix
b Equality constraint constant vector
adapt logical vector. Extra values to adapt, order must be the same as in $x$
nonneg logical vector of length(x). Determines which $x$-values have to obey nonnegativity constraints.
roundNearZeros Round near zero values of A before determining the sign?
tol tolerance used for zero-rounding.

## Details

There is some added flexibility. Users my define 'extra missings' by specifying the adapt vector.
By default it is assumed that all values must obey the nonnegativity constraint. However this can be determined by specifying nonneg.

## See Also

deduImpute, solSpace

## Examples

```
# a simple example"
E <- editmatrix(c(
        "x1 + x2 + x3 == xt",
        "x1 >= 0", "x2>=0","x3>=0", "xt>=0"))
x <- c(x1=10, x2=NA, x3=5,xt=15)
# with deductiveZeros we get:
( I <- deductiveZeros(E,x) )
x[I] <- 0
any(violatedEdits(E,x))
# This example is taken from De Waal et al (2011) (Examples 9.1-9.2)
E <- editmatrix(c(
        "x1 + x2 == x3",
        "x2 == x4",
        "x5 + x6 + x7 == x8",
        "x3 + x8 == x9",
        "x9 - x10 == x11",
        "x6 >= 0",
        "x7 >= 0"
))
x <- c(
        x1 = 145,
        x2 = NA,
        x3 = 155,
        x4 = NA,
        x5 = 86,
        x6 = NA,
        x7 = NA,
        x8 = 86,
        x9 = NA,
        x10 = 217,
        x11 = NA)
# determine zeros:
I <- deductiveZeros(E,x)
# impute:
x[I]<- 0
```


## Description

Based on observed values and edit rules, impute as many variables deductively as possible.
If $E$ is an editset, imputation based on numerical rules (if any) is performed, and imputations violating extra edits are reverted. Next, this procedure is repeated for pure categorical rules. The results are combined and returned in a deducorrect object.

For categorical data: The funcion deductiveLevels is used to derive deductive imputations for as many fields as possible
For numerical data: Given (equality) rules and a number of values to impute or adapt, in some cases unique solutions can be derived. This function uses solSpace and deductiveZeros (iteratively) to determine which values can be imputed deductively. Solutions causing new violations of (in)equality rules are rejected by default by testing if the observed values can lead to a feasible record. This may be switched off by passing checkFeasibility=FALSE. This may be desirable for performance reasons. If adapt was computed with an error localization algorithm, such as editrules: : localizeErrors, the feasibility check is also not nessecary.

## Usage

```
deduImpute(E, dat, adapt = NULL, ...)
## S3 method for class 'editset'
deduImpute(E, dat, adapt = NULL, ...)
## S3 method for class 'editarray'
deduImpute(E, dat, adapt = NULL, ...)
## S3 method for class 'editmatrix'
deduImpute(E, dat, adapt = NULL,
    tol = sqrt(.Machine$double.eps), round = TRUE, ...)
```


## Arguments

E
dat
adapt (optional) A boolean array of dim(dat), e.g. the result editrules::localizeErrors(E,dat). Column names must match those of dat.
... arguments to be passed to solSpace (numerical data) or deductiveLevels (categorical data)
tol tolerance to use in solSpace and in deductiveZeros
round should the result be rounded?

## Value

A deducorrect-object

## Note

When adapt is not NULL, values in dat where adapt==TRUE are replaced with NA. The output may therefore contain missings at positions that were previously filled (with wrong values, according to adapt).

## References

T. De Waal, J. Pannekoek and S. Scholtus (2011) Handbook of statistical data editing Chpt 9.2.19.2.2

## See Also

deductiveZeros, solSpace, deductiveLevels

## Examples

```
#############################################
# ImPUTATION OF NUMERIC DATA
#################################################
# These examples are taken from De Waal et al (2011) (Examples 9.1-9.2)
E <- editmatrix(c(
    "x1 + x2 == x3",
    "x2 == x4",
    "x5 + x6 + x7 == x8",
    "x3 + x8 == x9",
    "x9 - x10 == x11",
    "x6 >= 0",
    "x7 >= 0"
))
dat <- data.frame(
    x1=c(145,145),
    x2=c(NA,NA),
    x3=c(155,155),
    x4=c(NA,NA),
    x5=c(NA, 86),
    x6=c(NA,NA),
    x7=c(NA,NA),
    x8=c(86,86),
    x9=c(NA,NA),
    x10=c(217,217),
    x11=c(NA,NA)
)
```

```
dat
d <- deduImpute(E,dat)
d$corrected
d$status
d$corrections
```

```
#############################################
```

\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#

# IMPUTATION OF CATEGORICAL DATA

# IMPUTATION OF CATEGORICAL DATA

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# Here's an example from Katrika (2001) [but see De Waal et al (2011), ex. 9.3)]

E <- editarray(c(
"x1 \%in\% letters[1:4]",
"x2 \%in\% letters[1:3]",
"x3 \%in\% letters[1:3]",
"x4 \%in\% letters[1:2]",
"if (x2 == 'c' \& x3 != 'c' \& x4 == 'a' ) FALSE",
"if (x2 != 'a' \& x4 == 'b') FALSE",
"if (x1 != 'c' \& x2 != 'b' \& x3 != 'a') FALSE",
"if (x1 == 'c' \& x3 != 'a' \& x4 == 'a' ) FALSE"
))
dat <- data.frame(
x1 = c('c', NA ),
x2 = c('b',NA ),
x3 = c(NA ,NA ),
x4 = c(NA , 'b'),
stringsAsFactors=FALSE)
s <- deduImpute(E,dat)
s$corrected
s$status
s\$corrections
E <- editset(expression(
x + y == z,
x >= 0,
A %in% c('a','b'),
B %in% c('c','d'),
if ( A == 'a' ) B == 'b',
if ( B == 'b' ) x > 0
))
x <- data.frame(
x = NA,

```
```

        y = 1,
        z = 1,
        A = 'a',
        B = NA
    )

# deduImpute will impute x=0 and B='b',which violates the

# last edit. Hence, imputation will be reverted.

deduImpute(E,x)

```
getSignCorrection workhorse for correctSigns

\section*{Description}

Workhorse for correctSigns

\section*{Usage}
getSignCorrection(r, A1, C1, eps, A2, C2, epsvec, flip, swap, w, swapIsOneFlip, maxActions, maxCombinations)

\section*{Arguments}
\(r\)
A1 Equality matrix
C1 Constant vector for equalities
eps Tolerance for equality-checking
A2 Inequality matrix
C2 Constant vector for inequalities
epsvec Vector to check against. (.Machine.double.eps for > inequalities, otherwise 0.\()\)
flip indices in \(r\), where \(r\) may be sign-flipped
swap \(\quad n \times 2\) matrix with indices in \(r\), each row indicating a possible value swap.
w
swapIsOneFlip logical. If TRUE, weights are assigned to each flip or swap action. If FALSE weights are assigned to every changed variable.
maxActions Maximum number of flips+swaps to try.
maxCombinations
the maximum number of flip/swap combinations to try. See the description in correctSigns

\section*{Value}
a list containing
\begin{tabular}{ll} 
S & \(\mathrm{n} \times\) length \((r)\) array with corrected versions of \(r\) \\
weight & vector of length \(n\) with total weight for each solution \\
nFlip & number of sign flips for every solution \\
nSwap & number of value swaps for every solution
\end{tabular}

See Also
```

correctSigns

```
hellerTompkins

Determine if a matrix is totally unimodular using Heller and Tompkins criterium.

\section*{Description}

This function is deducorrect internal

\section*{Usage}
hellerTompkins(A)

\section*{Arguments}

A An object of class matrix in \(\{-1,0,1\}^{m \times n}\). Each column must have exactly 2 nonzero elements. (This is tested by isTotallyUnimodular).

\section*{Value}

TRUE if matrix is unimodular, otherwise FALSE

\section*{See Also}
```

isTotallyUnimodular

```
imputess Impute values from solution space

\section*{Description}

Given a record \(x\) with observerd \(x_{o b s}\) and missing values \(x_{m i s s}\) under linear equality constraints \(A x=b\). The function solSpace returns the solution space which can be written as \(x_{\text {miss }}=\) \(x_{0}+C z\), where \(x_{0}\) is are a constant vector (of dimension \(\mathrm{d}=\) length \(\left(x_{\text {miss }}\right)\) ) and \(C\) a constant matrix of dimension \(d \times d\).

\section*{Usage}
imputess(x, \(x 0, C, z=N U L L, ~ t o l=s q r t(. M a c h i n e \$ d o u b l e . e p s)) ~\)

\section*{Arguments}
x
\(x 0 \quad x 0\) outcome of solSpace
C C outcome of solSpace
z
tol tolerance used to check which rows of \(C\) equal zero.

\section*{Details}

If \(C\) has rows equal to zero, then those missing values may be imputed deductively. For the other missing values, some \(z\) must be chosen or another imputation method used.
The function imputess imputes missing values in a vector \(x\), based on the solution space and some chosen vector \(z\). If no \(z\) is passed as argument, only deductive imputations are performend (i.e. some missings may be left).
If \(C\) is a named matrix (as returned by solSpace), rows of \(x 0\) and \(C\) are matched by name to \(x\). Otherwise it is assumed that the missings in \(x\) occur in the order of the rows in \(C\) (which is also the case when x 0 and C are computed by solSpace).

\section*{Examples}
```

\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#

# IMPUTATION OF NUMERIC DATA

\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#

# These examples are taken from De Waal et al (2011) (Examples 9.1-9.2)

E <- editmatrix(c(
"x1 + x2 == x3",
"x2 == x4",
"x5 + x6 + x7 == x8",
"x3 + x8 == x9",
"x9 - x10 == x11",
"x6 >= 0",
"x7 >= 0"
))
dat <- data.frame(
x1=c(145,145),
x2=c(NA,NA),
x3=c(155,155),
x4=c(NA,NA),
x5=c(NA, 86),
x6=c(NA,NA),
x7=c(NA,NA),

```
```

    x8=c(86,86),
    x9=c(NA,NA),
    x10=c(217,217),
    x11=c(NA,NA)
    )
dat
d <- deduImpute(E,dat)
d$corrected
d$status
d\$corrections

```
```

\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#

# IMPUTATION OF CATEGORICAL DATA

\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#

```
\# Here's an example from Katrika (2001) [but see De Waal et al (2011), ex. 9.3)]
E <- editarray (c (
    "x1 \\%in\\% letters[1:4]",
    "x2 \\%in\\% letters[1:3]",
    "x3 \\%in\\% letters[1:3]",
    "x4 \\%in\\% letters[1:2]",
    "if (x2 == 'c' \& x3 != 'c' \& x4 == 'a' ) FALSE",
    "if (x2 != 'a' \& x4 == 'b') FALSE",
    "if (x1 != 'c' \& x2 != 'b' \& x3 != 'a') FALSE",
    "if (x1 == 'c' \& x3 != 'a' \& x4 == 'a' ) FALSE"
))
dat <- data.frame(
    \(\mathrm{x} 1=\mathrm{c}\left(\mathrm{c}^{\prime}, \mathrm{NA}\right)\),
    \(x 2=c(' b ', N A)\),
    x3 = c(NA , NA ),
    x4 = c(NA , 'b'),
    stringsAsFactors=FALSE)
s <- deduImpute(E, dat)
s\$corrected
s\$status
s\$corrections
E <- editset(expression(
    \(x+y==z\),
    \(x>=0\),
    A \%in\% c('a','b'),
    B \%in\% c('c','d'),
```

    if ( A == 'a' ) B == 'b',
    if ( B == 'b' ) x > 0
    ))
x <- data.frame(
x = NA,
y = 1,
z = 1,
A = 'a',
B = NA
)

# deduImpute will impute }\textrm{x}=0\mathrm{ and }\textrm{B}='\textrm{b}',\mathrm{ ,which violates the

# last edit. Hence, imputation will be reverted.

deduImpute(E,x)

```
    isTotallyUnimodular Test for total unimodularity of a matrix.

\section*{Description}

Check wether a matrix is totally unimodular.

\section*{Usage}
isTotallyUnimodular(A)

\section*{Arguments}

A An object of class matrix.

\section*{Details}

A matrix for which the determinant of every square submatrix equals \(-1,0\) or 1 is called totally unimodular. This function tests if a matrix with coefficients in \(\{-1,0,1\}\) is totally unimodular. It tries to reduce the matrix using the reduction method described in Scholtus (2008). Next, a test based on Heller and Tompkins (1956) or Raghavachari is performed.

\section*{Value}
logical

\section*{References}

Heller I and Tompkins CB (1956). An extension of a theorem of Danttzig's In kuhn HW and Tucker AW (eds.), pp. 247-254. Princeton University Press.
Raghavachari M (1976). A constructive method to recognize the total unimodularity of a matrix. _Zeitschrift fur operations research_, \({ }^{*} 20^{*}\), pp. 59-61.

Scholtus \(S\) (2008). Algorithms for correcting some obvious inconsistencies and rounding errors in business survey data. Technical Report 08015, Netherlands.

\section*{Examples}
```


# Totally unimodular matrix, reduces to nothing

A <- matrix(c(
1,1,0,0,0,
-1,0,0,1,0,
0,0,01,1,0,
0,0,0,-1,1),nrow=5)
isTotallyUnimodular(A)

# Totally unimodular matrix, by Heller-Tompson criterium

A <- matrix(c(
1,1,0,0,
0,0,1,1,
1,0,1,0,
0,1,0,1),nrow=4)
isTotallyUnimodular(A)

# Totally unimodular matrix, by Raghavachani recursive criterium

A <- matrix(c(
1,1,1,
1,1,0,
1,0,1))
isTotallyUnimodular(A)

```
newdeducorrect Generate an S3 deducorrect object

\section*{Description}

Generate an S3 deducorrect object

\section*{Usage}
newdeducorrect(corrected, corrections, status, Call = sys.call(-1))

\section*{Arguments}
corrected The corrected data.frame
corrections A data.frame listing old and new values for every row and variable where corrections were applied
status A data.frame with at least one status column.
Call Optionally, a call object.

\section*{Value}
an S3 object of class deducorrect

\section*{See Also}
```

    deducorrect-object
    ```
raghavachari \begin{tabular}{l} 
Determine if a matrix is unimodular using recursive Raghavachari cri- \\
terium
\end{tabular}

\section*{Description}

This function is deducorrect internal

\section*{Usage}
raghavachari(A)

\section*{Arguments}

A An object of class Matrix in \(\{-1,0,1\}^{m \times n}\).

\section*{Value}

TRUE or FALSE

See Also
isTotallyUnimodular
reduceMatrix Apply reduction method from Scholtus (2008)

\section*{Description}

Apply the reduction method in the appendix of Scholtus (2008) to a matrix. Let \(A\) with coefficients in \(\{-1,0,1\}\). If, after a possible permutation of columns it can be written in the form \(A=[B, C]\) where each column in \(B\) has at most 1 nonzero element, then \(A\) is totally unimodular if and only if \(C\) is totally unimodular. By transposition, a similar theorem holds for the rows of A. This function iteratively removes rows and columns with only 1 nonzero element from \(A\) and returns the reduced result.

\section*{Usage}
reduceMatrix(A)

\section*{Arguments}

A An object of class matrix in \(\{-1,0,1\}^{m \times n}\).

\section*{Value}

The reduction of A .

\section*{References}

Scholtus S (2008). Algorithms for correcting some obvious inconsistencies and rounding errors in business survey data. Technical Report 08015, Netherlands.

\section*{See Also}
isTotallyUnimodular

\section*{solSpace Solution space for missing values under equality constraints}

\section*{Description}

Solution space for missing values under equality constraints
solSpace method for editmatrix
This function finds the space of solutions for a numerical record \(x\) with missing values under linear constraints \(A x=b\). Write \(x=\left(x_{o b s}, x_{m i s s}\right)\). Then the solution space for \(x_{m i s s}\) is given by \(x_{0}+C z\), where \(x_{0}\) is a constant vector, \(C\) a constant matrix and \(z\) is any real vector of dimension ncol (C). This function computes \(x_{0}\) and \(C\).

\section*{Usage}
```

solSpace(E, x, ...)

## S3 method for class 'editmatrix'

solSpace(E, x, adapt = logical(length(x)),
checkFeasibility = TRUE, ...)
\#\# S3 method for class 'matrix'
solSpace(E, x, b, adapt = logical(length(x)),
tol = sqrt(.Machine\$double.eps), ...)

```

\section*{Arguments}

E
\(x \quad\) a named numeric vector.
... Extra parameters to pass to solSpace.matrix
adapt
and editmatrix or equality constraint matrix

A named logical vector with variables in the same order as in \(x\)
```

checkFeasibility

```

Check if the observed values can lead to a consistent record
b
Equality constraint constant vector
tol tolerance used to determine 0 -singular values when determining generalized inverse and to round coefficients of C to zero. See MASS: : ginv.

\section*{Details}

The user can specify extra fields to include in \(x_{\text {miss }}\) by specifying adapt. Also note that the method rests on the assumtion that all nonmissng values of \(x\) are correct.
The most timeconsuming step involves computing the generalized inverse of \(A_{\text {miss }}\) using MASS: :ginv (code copied from MASS to avoid dependency). See the package vignette and De Waal et al. (2011) for more details.

\section*{Value}

A list with elements \(x 0\) and \(C\) or NULL if the solution space is empty

\section*{References}
T. De Waal, J. Pannekoek and S. Scholtus (2011) Handbook of statistical data editing Chpt 9.2.1

Venables, W. N. \& Ripley, B. D. (2002) Modern Applied Statistics with S. Fourth Edition. Springer, New York. ISBN 0-387-95457-0

\section*{See Also}
```

deduImpute, deductiveZeros

```

\section*{Examples}
```


# This example is taken from De Waal et al (2011) (Examples 9.1-9.2)

E <- editmatrix(c(
"x1 + x2 == x3",
"x2 == x4",
"x5 + x6 + x7 == x8",
"x3 + x8 == x9",
"x9 - x10 == x11",
"x6 >= 0",
"x7 >= 0"
))
dat <- data.frame(
x1=c(145,145),
x2=c(NA,NA),
x3=c (155,155),
x4=c(NA,NA),
x5=c(NA, 86),
x6=c(NA,NA),

```
```

    x7=c(NA,NA),
    x8=c(86,86),
    x9=c(NA,NA),
    x10=c(217, 217),
    x11=c(NA,NA)
    )

# example with solSpace method for editmatrix

# example 9.1 of De Waal et al (2011).

x <-t(dat)[,1]
s <- solSpace(E,x)
S

# some values are uniquely determined and may be imputed directly:

imputess(x,s$x0,s$C)

# To impute everything, we choose z=1 (arbitrary)

z <- rep(1,sum(is.na(x)))
(y <- imputess(x,s$x0,s$C,z))

# did it work? (use a tolerance in checking to account for machine rounding)

# (FALSE means an edit is not violated)

any(violatedEdits(E,y,tol=1e-8))

# here's an example showing that solSpace only looks at missing values unless

# told otherwise.

Ey <- editmatrix(c(
"yt == y1 + y2 + y3",
"y4 == 0"))
y <- c(yt=10, y1=NA, y2=3, y3=7,y4=12)

# since solSpace by default checks the feasibility, we get no solution (since

# y4 violates the second edit)"

solSpace(Ey,y)

# If we ask solSpace not to check for feasibility, y4 is left alone (although

# the imputed answer is clearly wrong).

(s <- solSpace(Ey,y,checkFeasibility=FALSE))
imputess(y, s$x0, s$C)

# by setting 'adapt' we can include y4 in the imputation Since we know that

# with this adapt vector, imputation can be done consistently, we save some

# time by switching the feasibility check off.

(s <- solSpace(Ey, y, adapt=c(FALSE,FALSE,FALSE,FALSE,TRUE),
checkFeasibility=FALSE))
imputess(y,s$x0,s$C)

```

\section*{Description}

Create empty status vector

\section*{Usage}
status(n, ini = NA)

\section*{Arguments}
n
length of status vector
ini initial value, defaults to NA

\section*{Details}

Every function in deducorrect returns the status of every row after treatment. The status vector is an ordered factor with levels
\begin{tabular}{ll} 
invalid & record is invalid but could not be corrected \\
partial & record violates less edits then before entering the function \\
corrected & record satisfies all edit restrictions after correction \\
valid & record violates no edit restrictions
\end{tabular}
where invalid < partial < corrected < valid
This function is deducorrect internal.

\section*{Value}
an ordered factor with levels mentioned under details

\section*{Examples}
```


# create statusvector

status <- deducorrect:::status(5)
status[1:5] <- c("invalid",NA,"corrected","valid","partial")

# 

which(status < "valid")

```

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```
```

